Friday, 10 March 2023 OSCOAP is more efficient / weable in low-resource environment than DTLS I o T uses mostly UDP and CoAP

Transport layer application layer could roughly be HTTP adjusted for low-rower environment TCP is traditionally used for reliable connections UDP is connection less, and does not require maintaining connection (and resources required for this) delivery + duplicate protection not guaranteed delegated to application layer HTTP assumes TCP at tramport layer ... the border router generally has a propey translating between HTTP and COAP servery may also use COAP Co A V supports use as publish + subscribe rytocol Co AP response orderore similar to those of HTTP tokens are used to correlate responses with requests CoAP logically has two layers':

nemage layer

request/response layer Co AP header structure the nessage ID is welto detect displicates Confirmable messages regiure an ACK with the same message IDS non-confirmable messages do not (reset options are used when (non-) confirmable nessages are not supported) CoAP has caching features Co AP resource dissovery multicast address for list of available resources: GET /. well known/one CoAPhas 4 security modes O. no Rewrity (delegate to another layer) note that security at MAC layer allows intermediate modes to access content . Me - shored beg Me-programmed symmetric key 2. Taw public bay symmetric bey pre-programmed, no PKI 3. Certificates X.509 Certificates signed by CA, used for by eachange ECDSA for authentication ECDHE (with epheneral typ) for they agreement Co AP recommends the use of DTLS Co AP does not specify headers, setup of sewe layer, or how to use kys DTLS and TLS are similarly different as CoAP and HTTP datagram transport layer Dewrity DTLS provides end-to-end searcity 'at transport layer' ontop of UDP/TCP Toughly equivalent security guarantees, without order protection/non-replayability D TLS has coplicit sequence numbers DTLS (in particular handshoke) allows for re-ordering rackets & silently discording messages DTLS adas ACK messages DT LS merages carbe fragmented & re-assembled DTLS adds Hello Retry Request to avoid Mit M & Dos padet loss management is lased on a re-transmission timer initially 9 s, retransmit upon sering, the double timer "levery layer:

- sequence number allow re-ordering, decrypting individual reases, protecting against replay

- epoch allows identification of bey charges only block ciphers, no stream cirkery ore allowed hardshake pertocal same as TLS VI.3, but adjusted for UDP DTLS handstake dienthelle e.g. supported cipher suites Relloretry represt + cookie - must be re-set with askie; used to prevet storm of grouped packets to you fed clienthello IP selected in her mite nersyes are grouped to conserve energy DTLS is the main protocol inits niche, but it's not the most efficient object searcity for CoAP 05 CoAP multiple clients correct to one server encryption signing of Co AP message using me - Shored by material authorization server parts beyon on client and access server and access server and access server (Ner client) on server incertiging they have the server to be a Client asks resource Derver returns encrypted (and signed) resource bey is computed from access secret, message ID and sender ID advantage: dels not require hardsbake & record protocols Co AP flocking e.g. (selectively) blocking requests/regrosson
mitigation through e.g. energetim (does not always work), attack detection through confirmable messages, \$\pm\0505 C. AP respect delay delay reaching destinations mitigating through IDS, shortening transmission window (but gives reduced robustness to network fluctuations) Co AP relay attacks records packets the relay then to other devices Cor by with remote unlock mitigation: neasure RTT (round-trip time), use (rigned) 6PS location data MQTT message quering telemetry transport rublish mbscribe messaging transport protocol broker ('server): hosts list of topics + subscribers, redirects updates on torics to subscribers nodes: subwrife to topics of interest, "eceives updates on topics publisher publishes updates on topic, Canalso be a subscriber MQTT runs over TCP wit UDP -> nessages may not receive out of order in MRTT MOTT-SN rung over UDP and has smaller headers and the concept of sleeping clients ensor notwork MQT T security features authertication - frust usernames password, later challenge-response authorization - we toke / usernane to beck permissions integrity - Rules or suggestion to use VPN confishentiality

many possible anomalies exist in MQTT

Lecture 8